

Cloud Computing

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Key Reference:

Prof. Jong-Moon Chung's Lecture Notes at Yonsei University



Cloud Computing

- Cloud Introduction
- Cloud Service Model
- Big Data
- Hadoop
- MapReduce
- HDFS (Hadoop Distributed File System)







- Hadoop
- Hadoop is a Reliable Shared Storage and Analysis System
- Hadoop = HDFS + MapReduce + α
 - HDFS provides Data Storage
 - **HDFS**: Hadoop Distributed FileSystem
 - MapReduce provides Data Analysis
 - MapReduce = Map Function + Reduce Function



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- Scaling Out
- Scaling out is done by the DFS (Distributed FileSystem), where the data is divided and stored in distributed computers & servers
- Hadoop uses HDFS to move the MapReduce computation to several distributed computing machines that will process a part of the divided data assigned





Jobs

- MapReduce job is a unit of work that needs to be executed
- Job types: Data input, MapReduce program, Configuration Information, etc.
- Job is executed by dividing it into one of two types of tasks
 - Map Task
 - Reduce Task

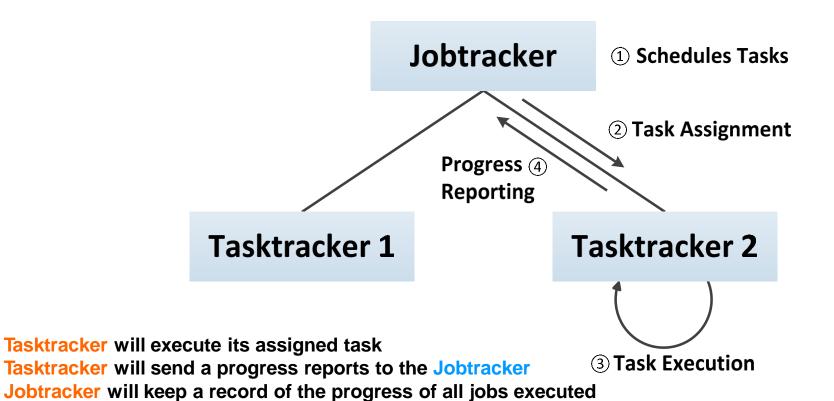


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- Node types for Job execution
 - Job execution is controlled by 2 types of nodes
 - Jobtracker
 - Tasktracker
 - Jobtracker coordinates all jobs
 - Jobtracker schedules all tasks and assigns the tasks to tasktrackers









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- Data flow
 - Hadoop divides the input into input splits (or splits) suitable for the MapReduce job
 - Split has a fixed-size
 - Split size is commonly matched to the size of a HDFS block (64 MB) for maximum processing efficiency



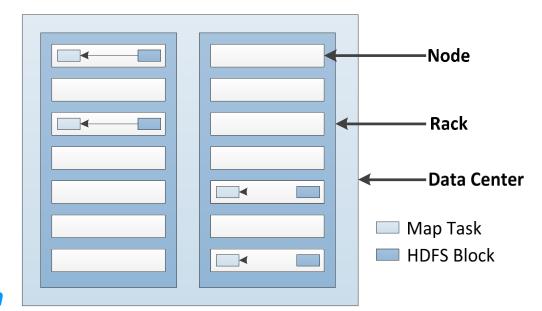


- Data flow
 - Map Task is created for each split
 - Map Task executes the map function for all records within the split
 - Hadoop commonly executes the Map Task on the node where the input data resides



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Data flow



- Data-Local Map Task
- Data locality optimization

does not need to use the cluster network

Data-local flow process shows why the

Optimal Split Size = 64 MB HDFS Block Size



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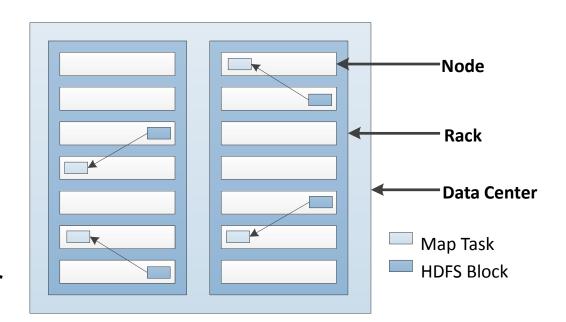
Data flow



 A node hosting the HDFS block replicas for a map task's input split

could be running other map tasks

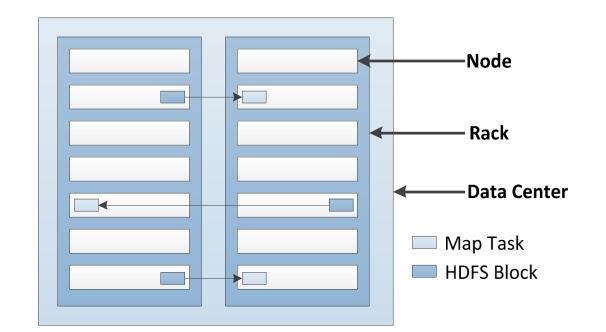
 Job Scheduler will look for a free map slot on a node in the same rack as one of the blocks





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Data flow



- Off-Rack Map Task
- Needed when the Job Scheduler

cannot perform data-local or rack-local map tasks

Uses inter-rack network transfer



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Map

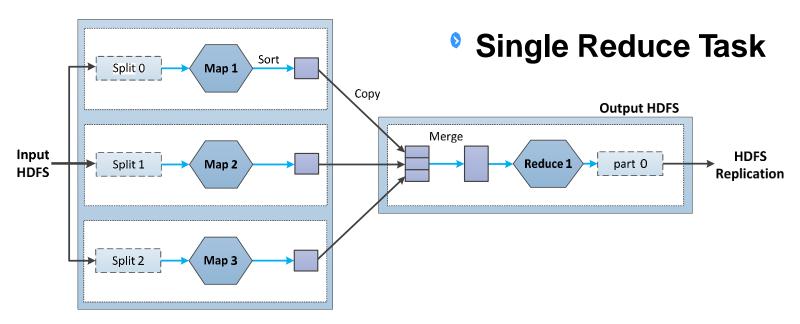
- Map task will write its output to the local disk
- Map task output is not the final output, it is only the intermediate output

Reduce

- Map task output is processed by Reduce Tasks to produce the final output
- Reduce Task output is stored in HDFS
 - For a completed job, the Map Task output can be discarded



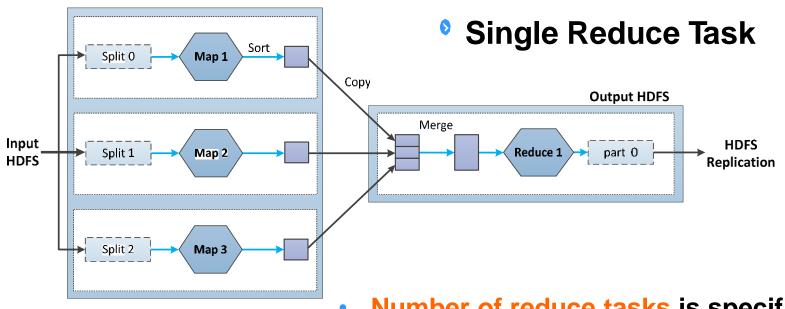




- Node includes Split, Map, Sort, and Output unit
- Light blue arrows show data transfers in a node
- Black arrows show data transfers between nodes







 Number of reduce tasks is specified independently, and is not based on the size of the input



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- Combiner Function
 - User specified function to run on the Map output
 - → Forms the input to the Reduce function
 - Specifically designed to minimize the data transferred between Map Tasks and Reduce Tasks
 - Solves the problem of limited network speed on the cluster and helps to reduce the time in completing MapReduce jobs

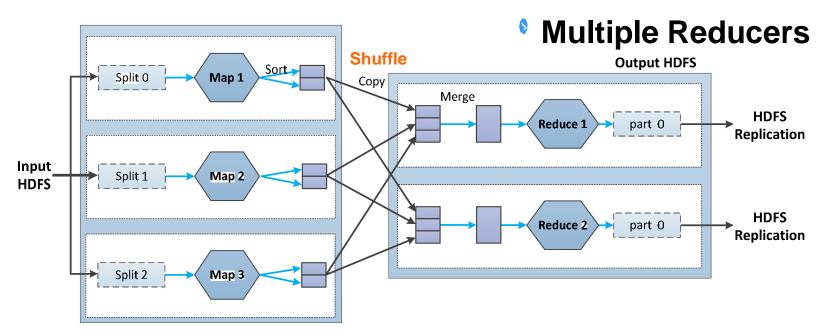




- Multiple Reducer
 - Map tasks partition their output, each creating one partition for each reduce task
 - Each partition may use many keys and key associated values
 - All records for a key are kept in a single partition







 Shuffle process is used in the data flow between the Map tasks and Reduce tasks

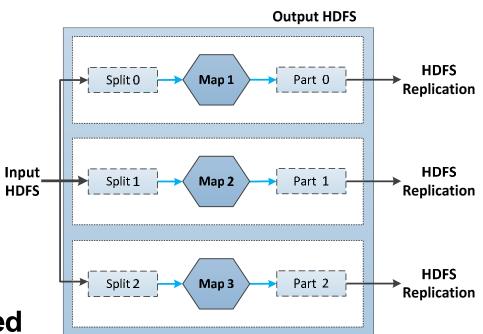


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Zero Reducer

 Zero reducer uses no shuffle process

 Applied when all of the processing can be carried out in parallel Map tasks









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 - **HDFS**: Hadoop Distributed FileSystem
 - MapReduce provides Data Analysis
 - MapReduce = Map Function + Reduce





- HDFS: Hadoop Distributed FileSystem
 - DFS (Distributed FileSystem) is designed for storage management of a network of computers
 - HDFS is optimized to store large terabyte size files with streaming data access patterns



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- HDFS: Hadoop Distributed FileSystem
 - HDFS was designed to be optimal in performance for a WORM (Write Once, Read Many times) pattern
 - HDFS is designed to run on clusters of general computers
 & servers from multiple vendors





- HDFS Characteristics
- HDFS is optimized for large scale and high throughput data processing
- HDFS does not perform well in supporting applications that require minimum delay (e.g., tens of milliseconds range)



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- Blocks
 - Files in HDFS are divided into block size chunks
 - **→** 64 Megabyte default block size
 - Block is the minimum size of data that it can read or write
 - Blocks simplifies the storage and replication process
 - → Provides fault tolerance & processing speed enhancement for larger files





- HDFS
 - HDFS clusters use 2 types of nodes
 - Namenode (master node)
 - Datanode (worker node)





- Namenode
 - Manages the filesystem namespace
 - Namenode keeps track of the datanodes that have blocks of a distributed file assigned
 - Maintains the filesystem tree and the metadata for all the files and directories in the tree
 - Stores on the local disk using 2 file forms
 - Namespace Image
 - Edit Log





- Namenode
 - Namenode holds the filesystem metadata in its memory
 - Namenode's memory size determines the limit to the number of files in a filesystem
 - But then, what is Metadata?



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Metadata

- Traditional concept of the library card catalogs
- Categorizes and describes the contents and context of the data files
- Maximizes the usefulness of the original data file by making it easy to find and use





- Metadata Types
 - Structural Metadata
 - Focuses on the data structure's design and specification
 - Descriptive Metadata
 - Focuses on the individual instances of application data or the data content



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Datanodes

- Workhorse of the filesystem
- Store and retrieve blocks when requested by the client or the namenode
- Periodically reports back to the namenode with lists of blocks that were stored



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Client Access

- Client can access the filesystem (on behalf of the user) by communicating with the namenode and datanodes
- Client can use a filesystem interface (similar to a POSIX (Portable Operating System Interface)) so the user code does not need to know about the namenode and datanodes to function properly



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- Namenode Failure
- Namenode keeps track of the datanodes that have blocks of a distributed file assigned
 - → Without the namenode, the filesystem cannot be used
- If the computer running the namenode malfunctions then reconstruction of the files (from the blocks on the datanodes) would not be possible
 - → Files on the filesystem would be lost





Namenode Failure Resilience

- Namenode failure prevention schemes
 - 1. Namenode File Backup
 - 2. Secondary Namenode





- Namenode File Backup
 - Back up the namenode files that form the persistent state of the filesystem's metadata
 - Configure the namenode to write its persistent state to multiple filesystems
 - → Synchronous and atomic backup
 - Common backup configuration
 - Copy to Local Disk and Remote FileSystem



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- Secondary Namenode
- Secondary namenode does not act the same way as the namenode
- Secondary namenode periodically merges the namespace image with the edit log to prevent the edit log from becoming too large
- Secondary namenode usually runs on a separate computer to perform the merge process because this requires significant processing capability and memory



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Hadoop 2.x Release Series HDFS Reliability Enhancements

- HDFS Federation
- HDFS HA (High-Availability)





- HDFS Federation
 - Allows a cluster to scale by adding namenodes
 - Each namenode manages a namespace volume and a block pool
 - Namespace volume is made up of the metadata for the namespace
 - Block pool contains all the blocks for the files in the namespace





HDFS Federation

- Namespace volumes are all independent
 - Namenodes do not communicate with each other
 - Failure of a namenode is also independent to other namenodes
 - A namenode failure does not influence the availability of another namenode's namespace



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- HDFS High-Availability
 - Pair of namenodes (Primary & Standby) are set to be in Active-Standby configuration
 - Secondary namenode stores the latest edit log entries and an up-to-date block mapping
 - When the primary namenode fails, the standby namenode takes over serving client requests



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- HDFS High-Availability
 - Although the active-standby namenode can takeover operation quickly (e.g., few tens of seconds), to avoid unnecessary namenode switching, standby namenode activation will be executed after a sufficient observation period (e.g., approximately a minute or a few minutes)



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