Three useful categories

Learning a programming language involves:

Syntax: The grammar rules defining a program fragment.

Semantics: The meaning of various programming fragments.

Pragmatics: How to effectively use language features, libraries, IDEs,

• • •

All three of these are important in how easy it is to easily write high-quality software.

For all categories, consider the principle of least surprise.

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• variable:

type:

• expression:

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- expression: a piece of syntax which evaluates to some particular value.
 - E.g. 3+4*5 or sqrt(16).

Some vocabulary (cont.)

• *literal*: a piece of syntax which evaluates immediately to a particular value.

E.g. Java 37 or 045 are both literals representing the value 37, which is of *type* int. And 37., 37d, 37e0 are each literal double s.

(We will often conflate a literal with the value it represents, and only say "literal" when we're emphasizing that we're dealing with syntax.)

Literals occur in the source-code text, and can be processed at compile-time. In Java, string literals are "interned": If the same string-literal occurs twice, the the compiler is smart enough to only make one object(*), and use the same reference in both places.

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typing: when?

statically-typed: At compile-time, the types of all declared names are known.

Can be provided by programmer and checked by type-system, or inferred by the language (ML, Haskell). (C# allows simple var n = 5; and infers $n \in int$).

dynamically-typed: Language knows the type of every value.

But a variable might hold values of different types, over its lifetime. php, javascript, racket. Each value may include some extra bits, indicating its type.

typing: other approaches

duck typing: Care about an object having a field/method, not any inheritance.

E.g. javascript

untyped:

E.g. assembly

type-safe: Any type error is caught (either dynamically or statically).

Note that C is not type-safe, due to casting. Java's casting is type-safe (since a bad cast will fail at run-time).

typing: strong/weak/non

These terms are often used in different ways:

strongly typed: no/few implicit type conversions, or statically typed

weakly typed: many implicit type conversions, or dynamically typed

Consider Java Math.sqrt(16), and Java vs php "50" + 60.

Compiling vs Interpreting

A compiler is a function

```
compile : source-code → machine-code
```

The resulting machine-code, when executed, runs the program which produces a resulting value.

An interpreter is a function

```
eval : expr \rightarrow value
```

which evaluates an expression, producing a resulting value.

Compiling vs Interpreting (cont.)

- Running interpreted code, you are running the interpreter, which is looking at the source-expression as if it were data.
- Running compiled code, you are running the program directly.
- Compiled code: faster, but platform-specific.

The distinction is practical, but not fundamental: You can even claim that CPUs are simply interpreters: they read compiled-code as data (from memory), and update their internal state accordingly.

• A compromise: compile to byte code (or, javascript), and run an interpreter for that byte code. A speed/platform-dependence trade-off.